

Covert Channels Through Branch Predictors: a Feasibility Study

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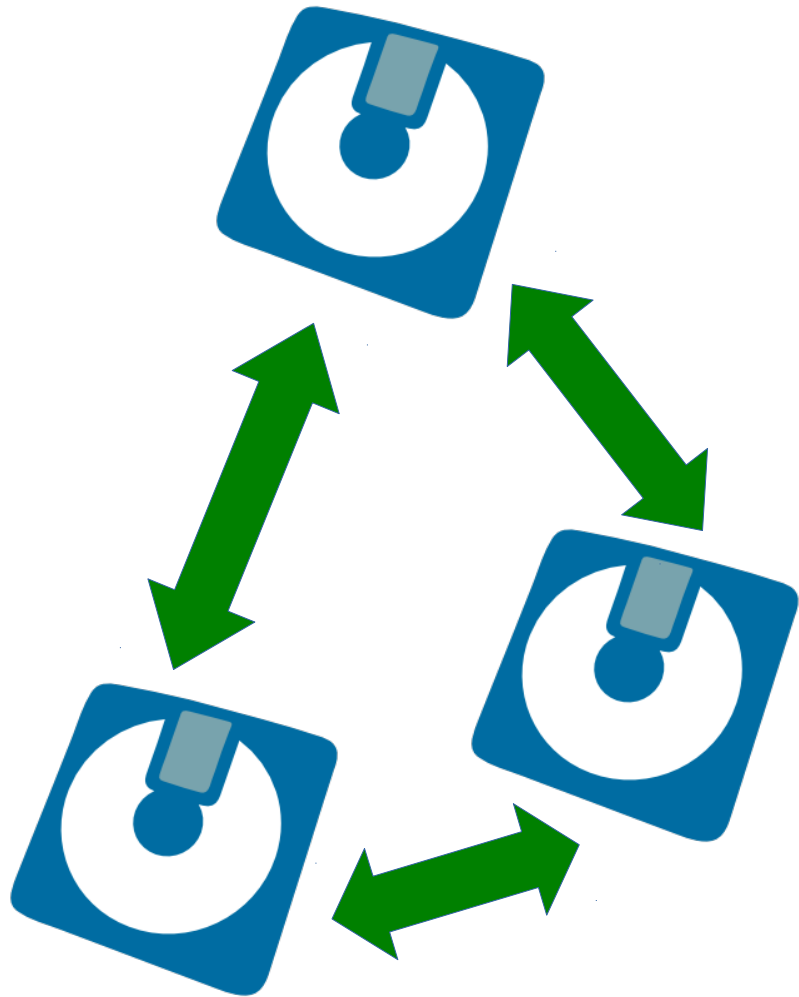
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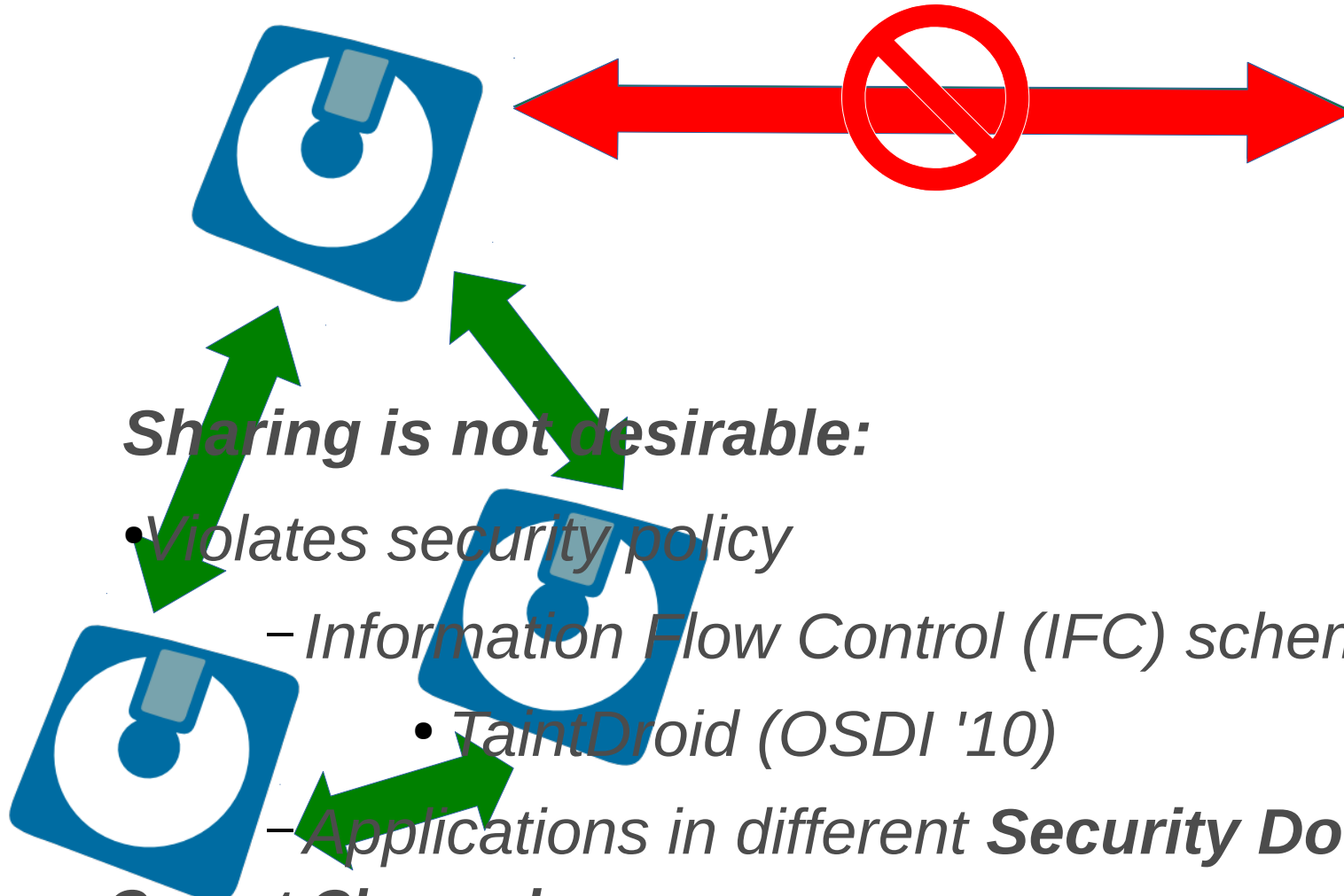


- *What is covert channel? How can it be used?*
- *Architectural covert channels*
- *Branch Predictor in CPU*
- *Constructing covert channel through Branch Predictor*
- *Results*
- *Optimizations*
- *Conclusion*



Data Sharing Mechanisms:

- *Shared Memory*
- *Networking*
- *File system*
- *IPC*
- *etc...*



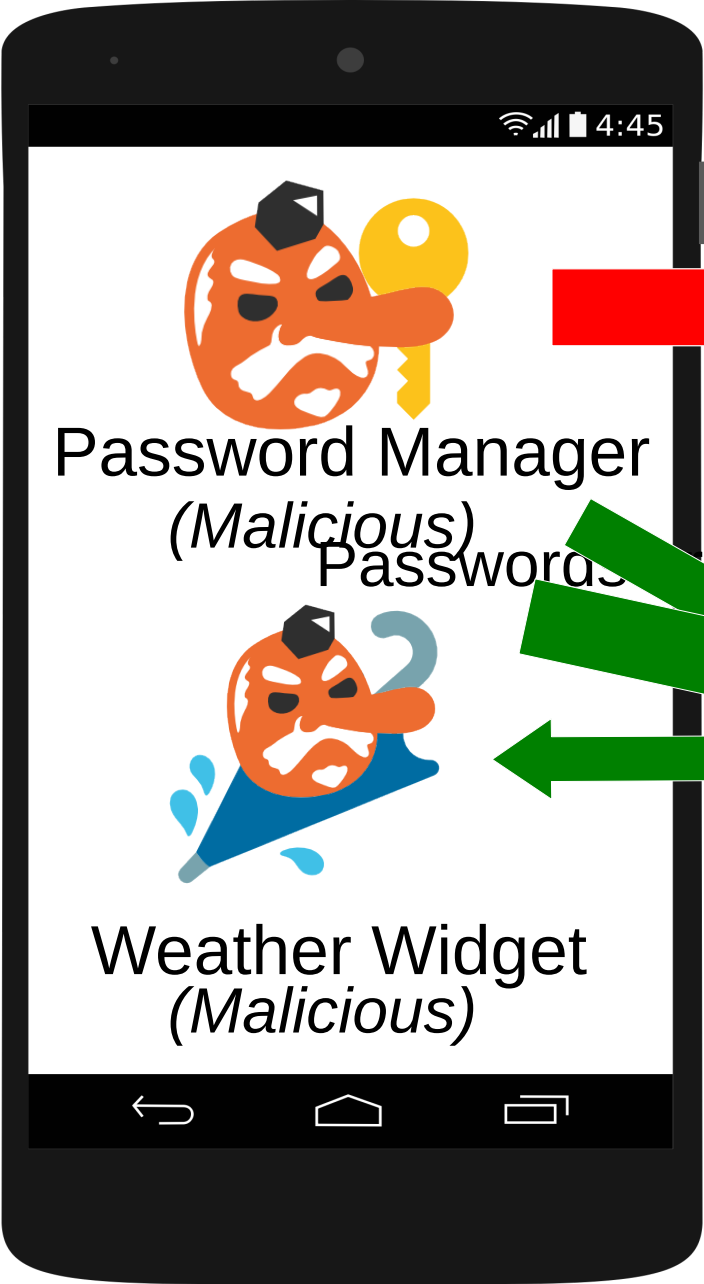
Sharing is not desirable:

- *Violates security policy*
 - *Information Flow Control (IFC) schemes*
 - *TaintDroid (OSDI '10)*
 - *Applications in different **Security Domains***

Covert Channels:

Transfer information through channels not intended for information transfer

Covert Channel Usage Example



Application is isolated
or password data is tainted

Are passwords secure?

No IPC (#&)
No File System Access
No Network Access



Passwords are secure as long as OS enforces security policies

Network Access

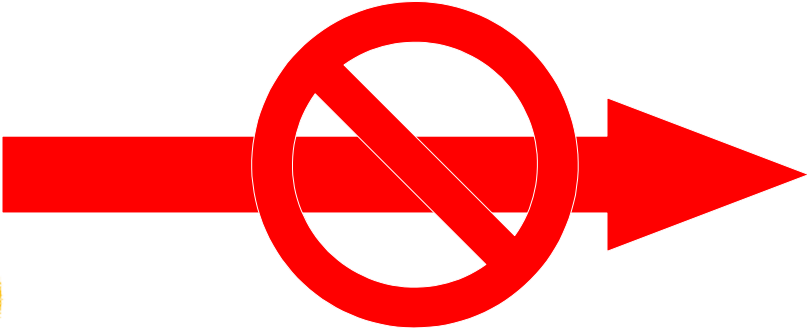
Covert Channel Usage Example

Shared resources can be used to construct covert channels:

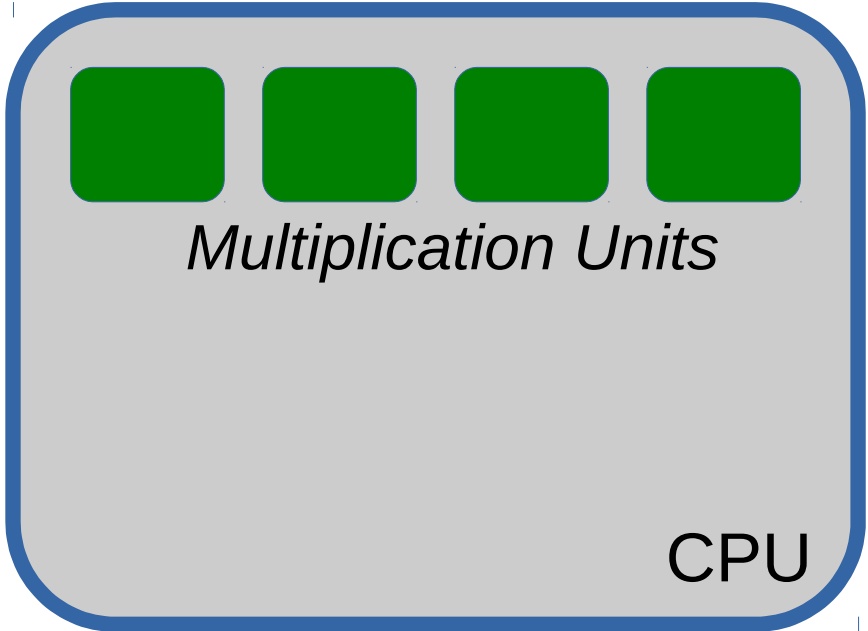
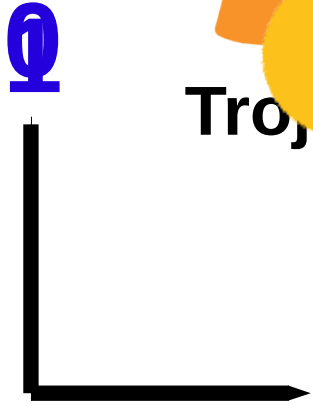
- *OS resources*
 - *File descriptors, free space, etc.*
- *Network latencies*
- *Power and Thermal effects*
- ***Architectural resources***
 - *Caches*
 - *Functional units*
 - ***Branch predictor***
 - *Other resources*
- *Others*

Contention-based Architectural Covert Channels

```
mul %r1,%r1  
mul %r1,%r1  
mul %r1,%r1  
mul %r1,%r1
```



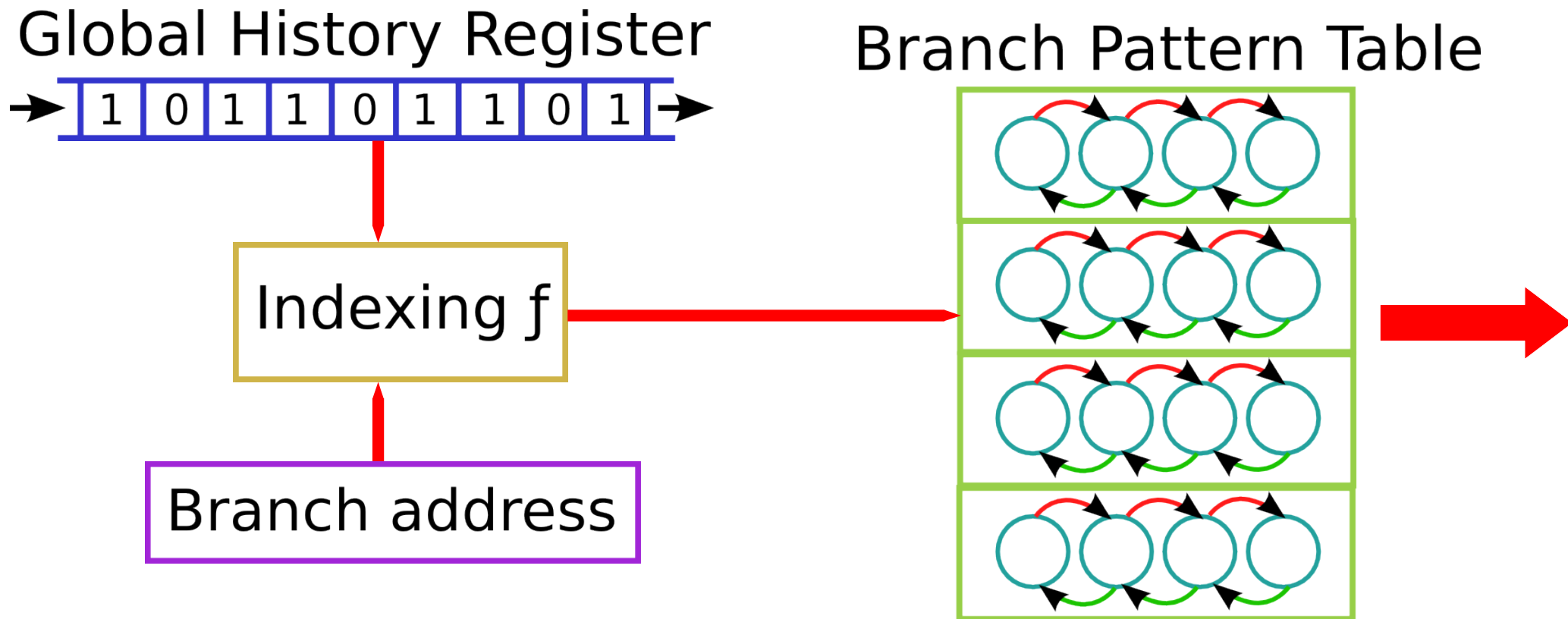
```
mul %r1,%r1  
mul %r1,%r1  
mul %r1,%r1  
mul %r1,%r1
```



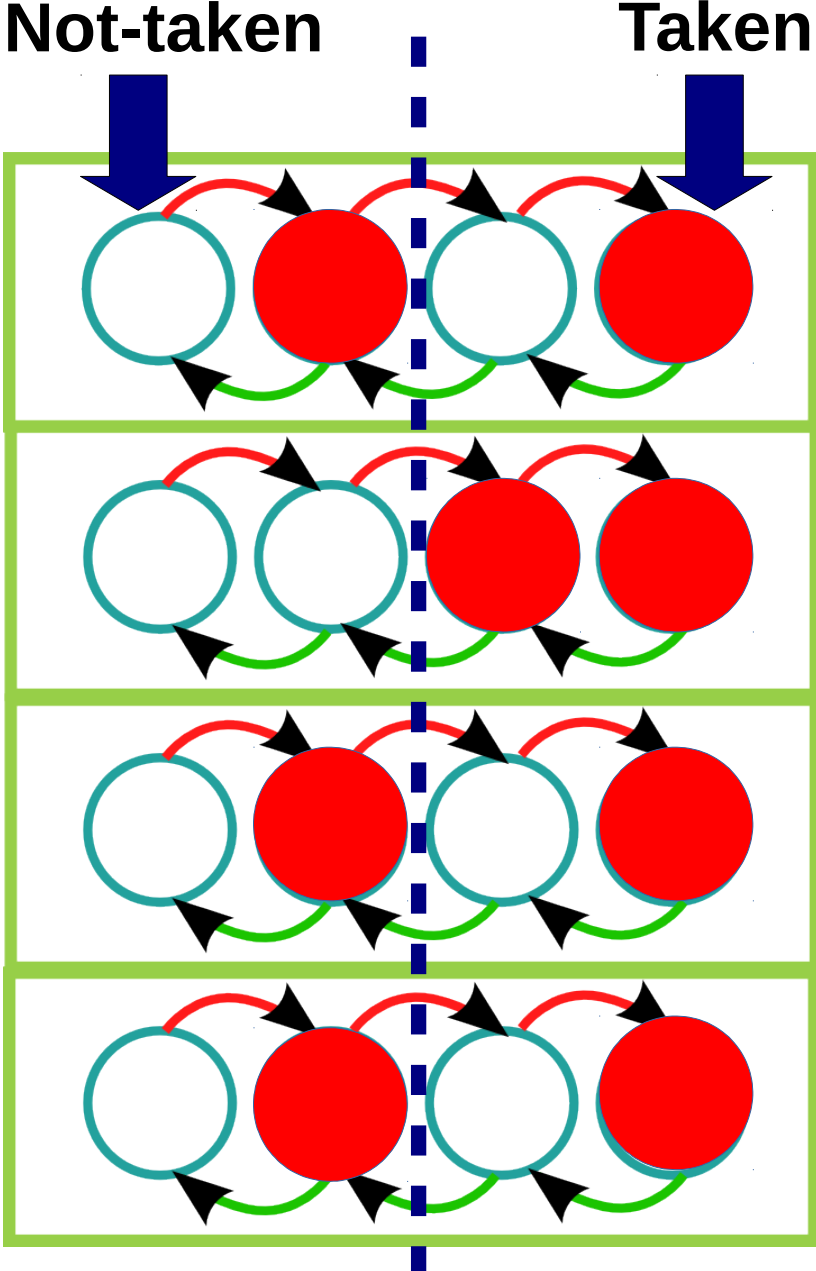
Properties that make covert channel possible:

- *During execution branch predictor accumulates state*
- *BP is shared among all processes on core*
- *BP is not flushed on context switches*
- *Parallel threads in SMT share same BP*
- *Branch mispredictions have high cost*

Branch Predictor Operation



Constructing Covert Channel



```
while(1)
  if (time(0)%2){
    taken();
  }else{
    nottaken();
  }
}
```



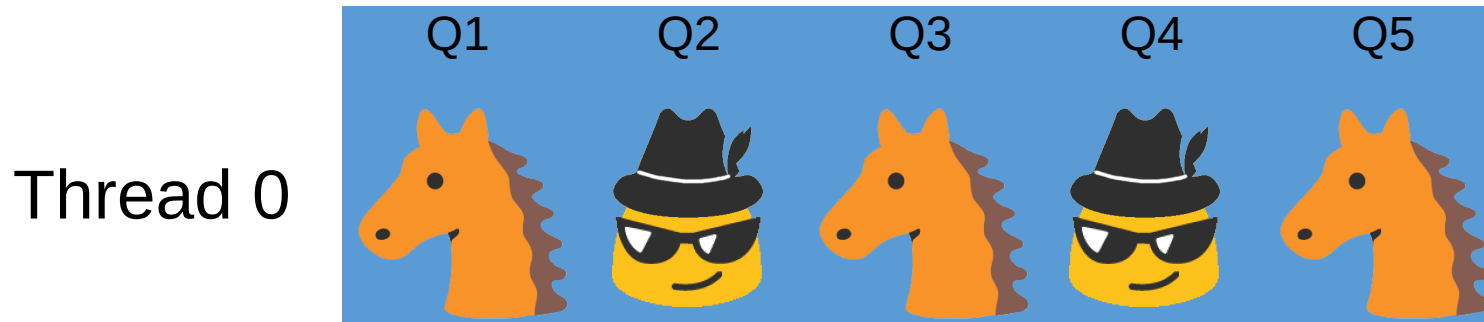
```
for (int i=0; i <
MAX_PROBES;i++){
  usleep(SLEEP_T);
  start_t=rdtsc();
  taken();
  end_t=rdtsc();
}
```



Scheduling Trojan and Spy

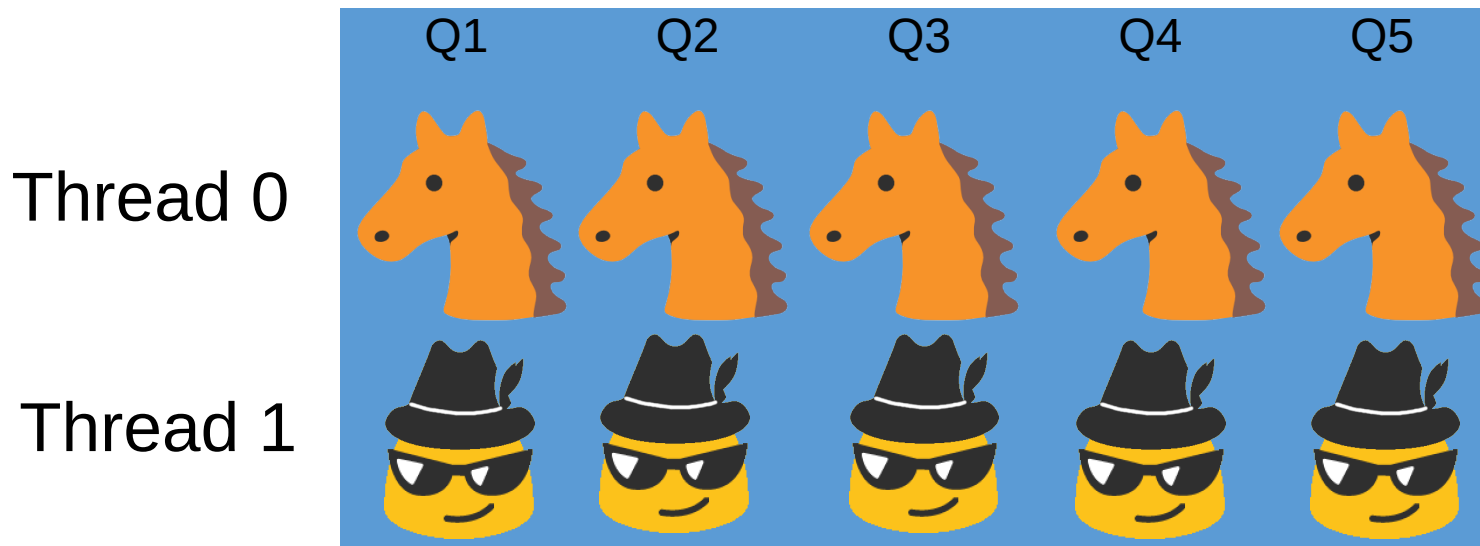
Single Threaded Scenario:

Setting affinity mask to run ~~Trojan and Spy~~ on the same core



Multi Threaded Scenario (SMT):

Setting affinity mask to run ~~Trojan and Spy~~ on parallel virtual cores



```
br_taken:
push  %rbp
movl  $0x1,-0x8(%rbp)
cmpl  $0x0,-0x8(%rbp)
jne   .L2
nop
nop
.L2:
cmpl  $0x0,-0x8(%rbp)
jne   .L3
nop
nop
nop
.L3:
.....
```



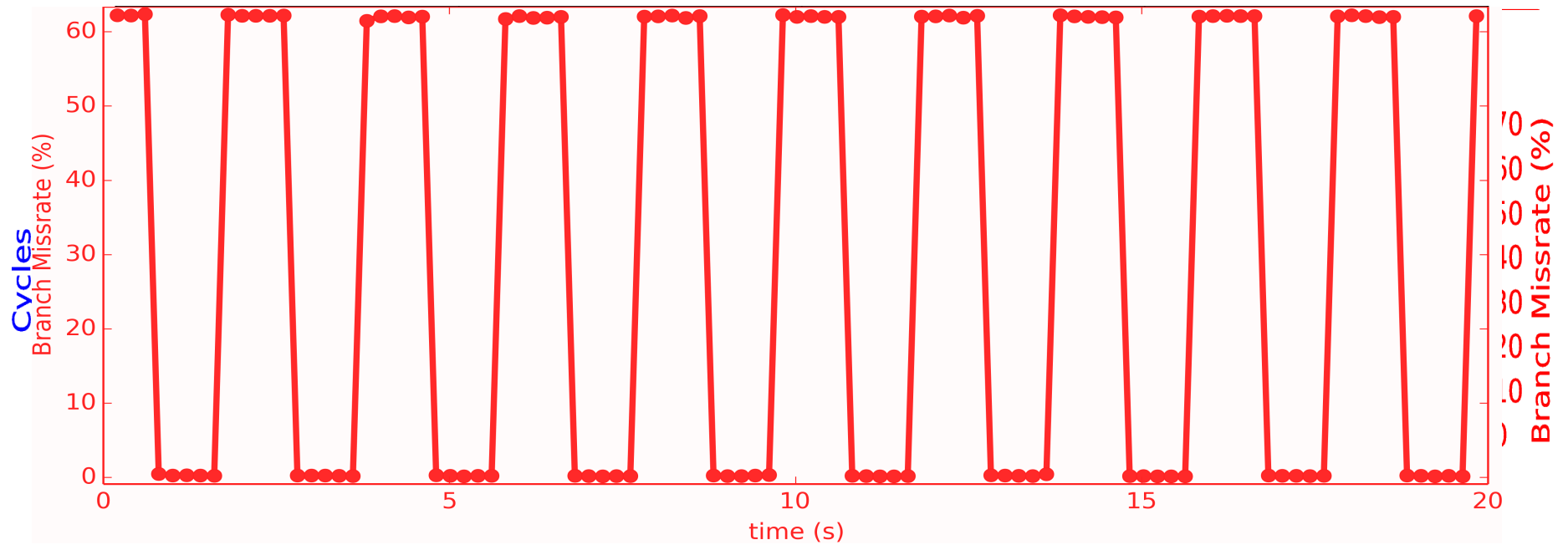
Notes:

- Trojan and Spy execute large number of branch instructions
 - 500K for Trojan
 - 30K for Spy
- Trojan and Spy execute the similar code
- Nop instructions to randomize layout

- Real hardware
- Intel Core i7-4800MQ CPU (Haswell)
- Ubuntu 14.04.2 generic Linux kernel version 3.16.0-31

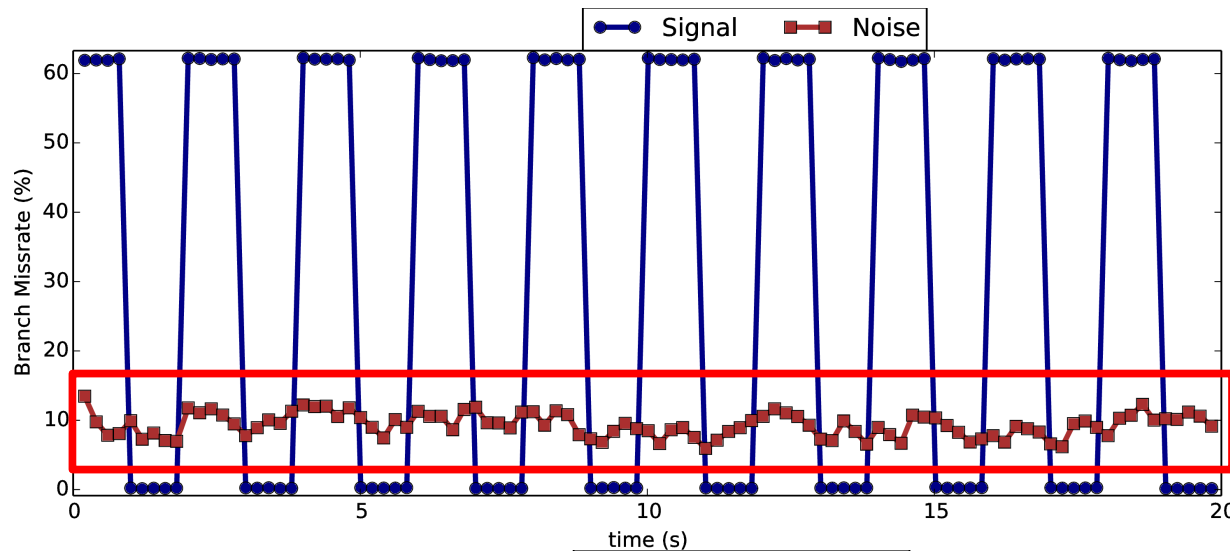
Results: Single Thread

Single Thread, no interference:

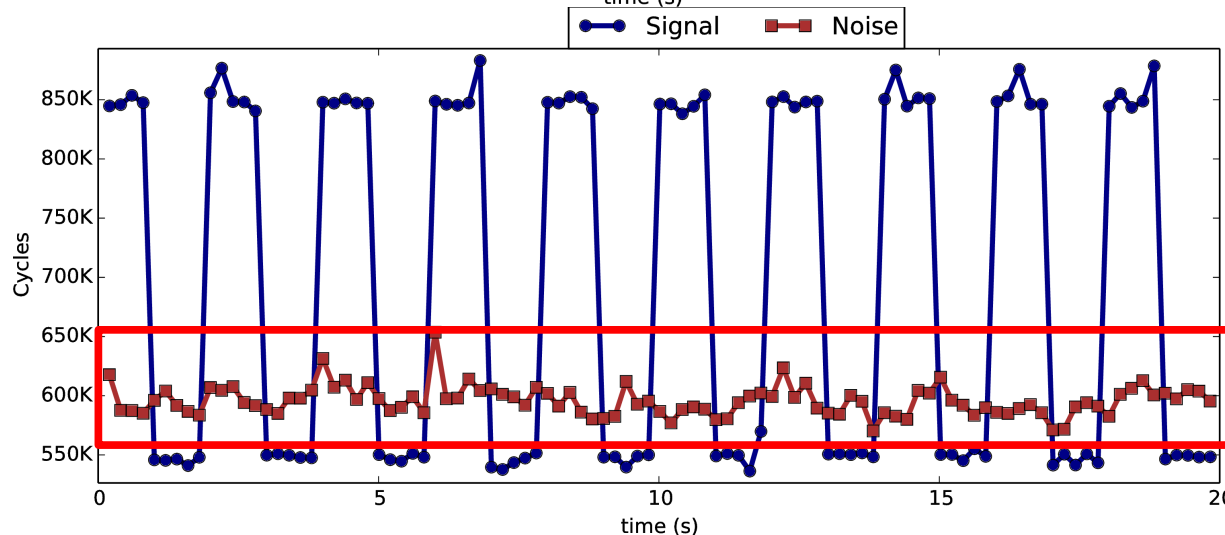


Signal vs Noise

Single Thread, no interference, *cpuburn* as noise:

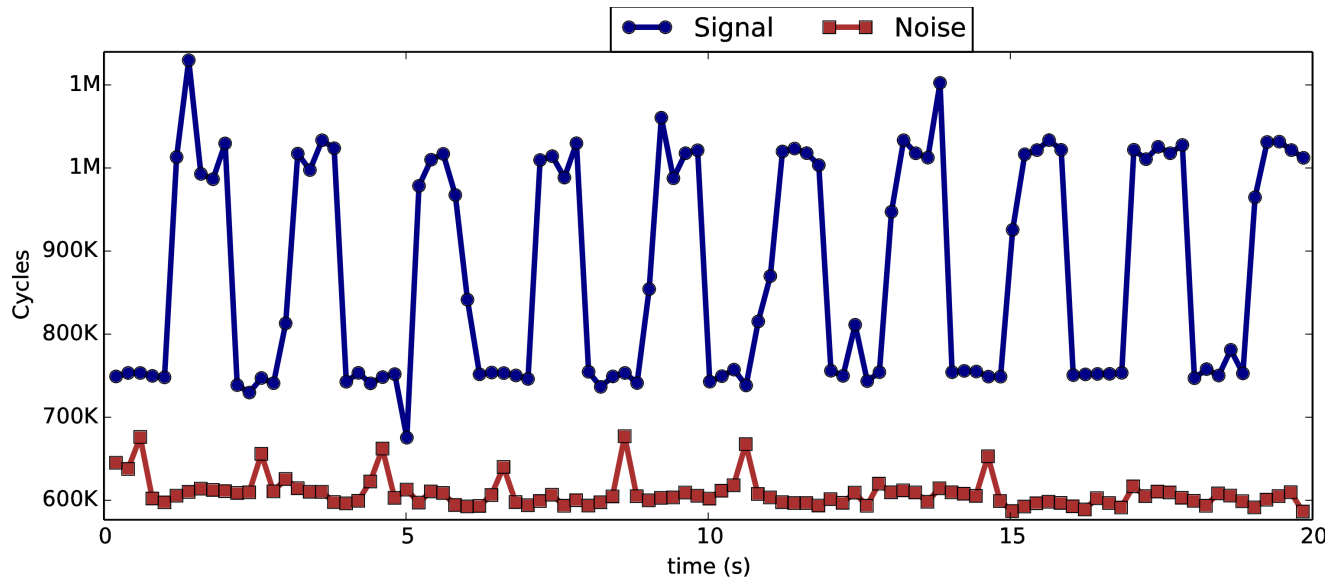
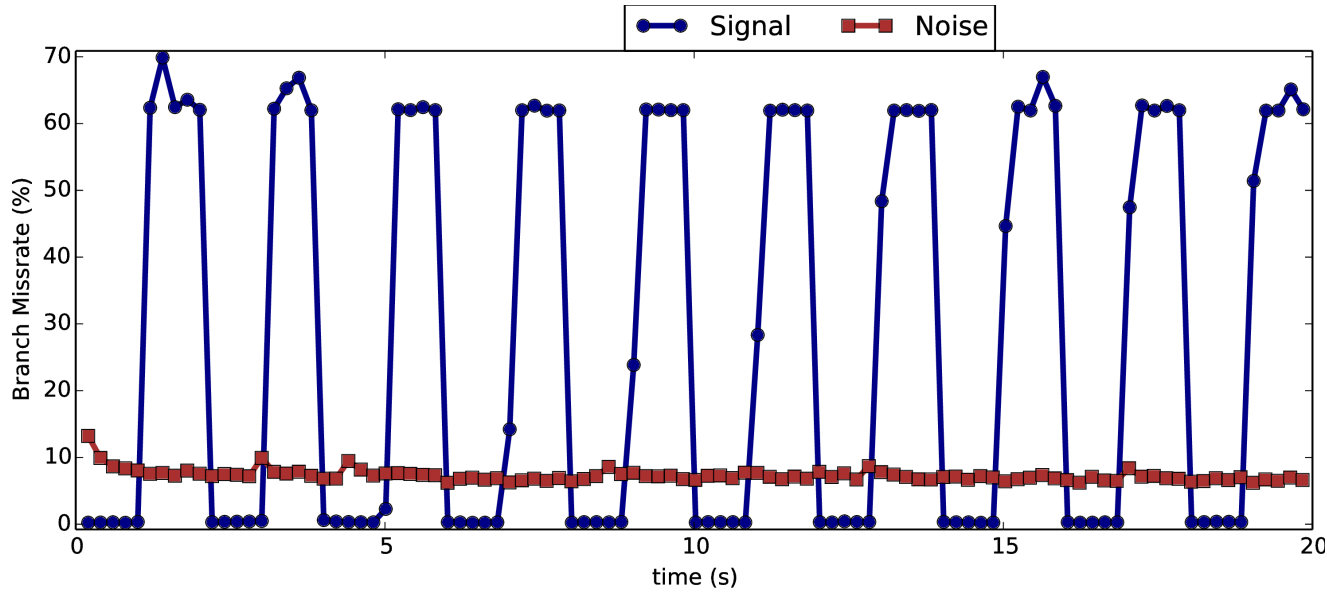


- High SNR
- Stable signal
- Possibly high capacity
- Asynchronous signal



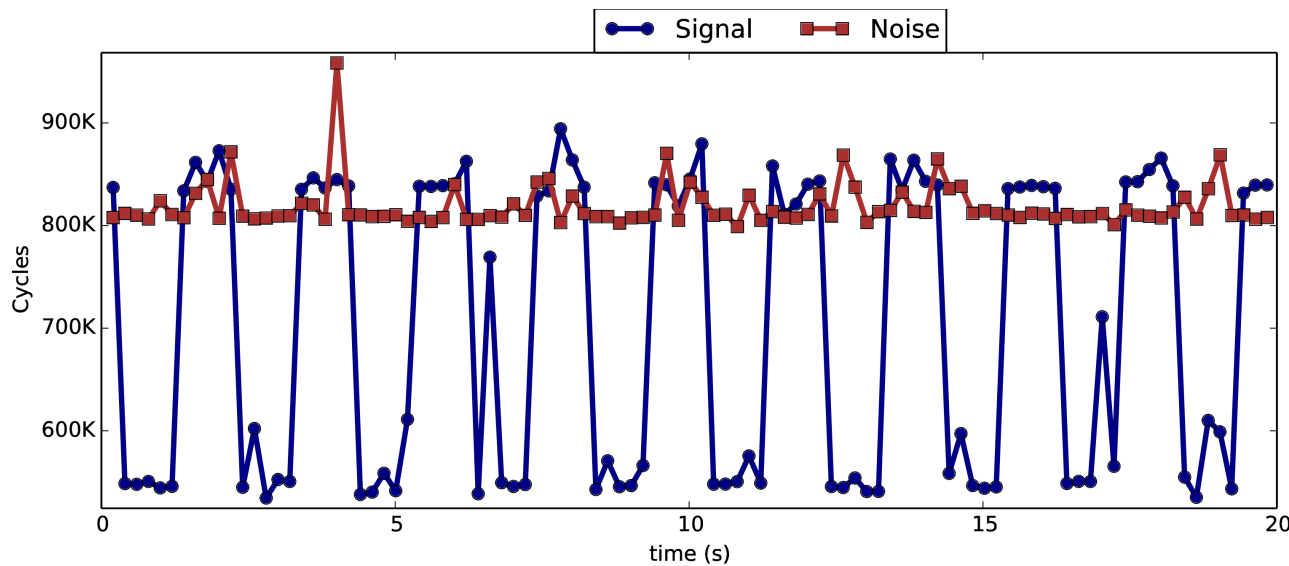
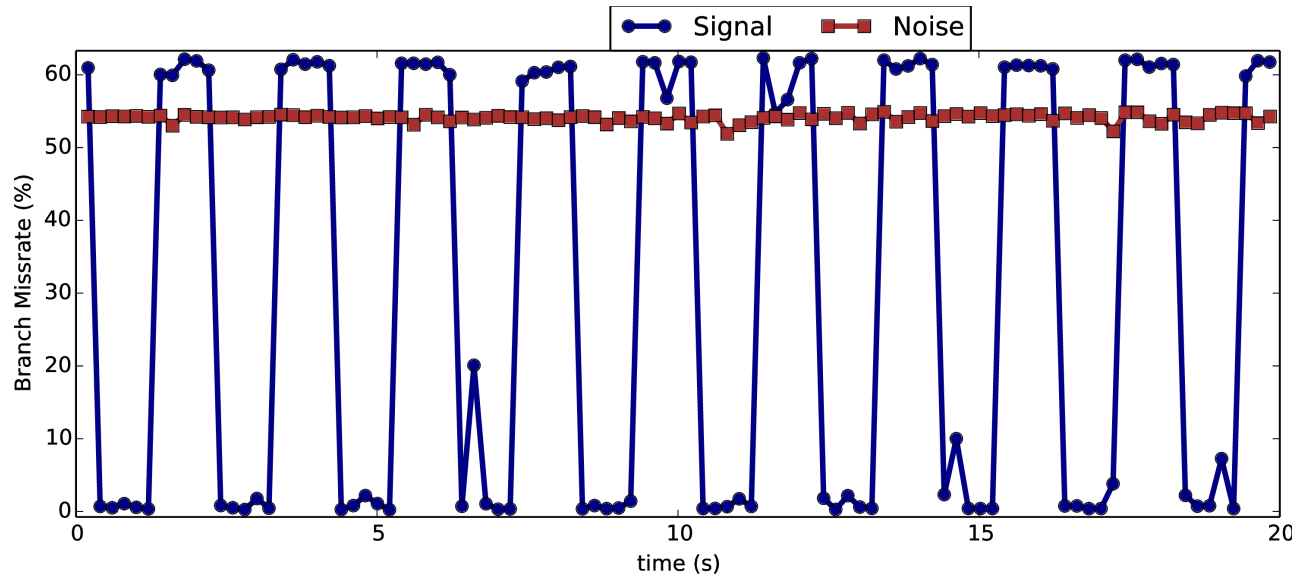
SMT Results

Multi Thread, no interference, *cpuburn* as noise:



Realistic example

Single Thread, Browser with YouTube as noise/interference



- *What we present is only a **prototype***
- *Can transmit multiple bits at once*
- *Can make scheduling faster*
- *For more accuracy, can reverse-engineer BP to find optimal number of branch instructions/nops*

- *Branch predictor can be used as an effective covert channel media*
- *Covert channel through BP has desirable properties*
- *It should be considered when implementing secure systems free of covert and side channels*

Thank you!