

Using Scan Side Channel to Detect IP Theft

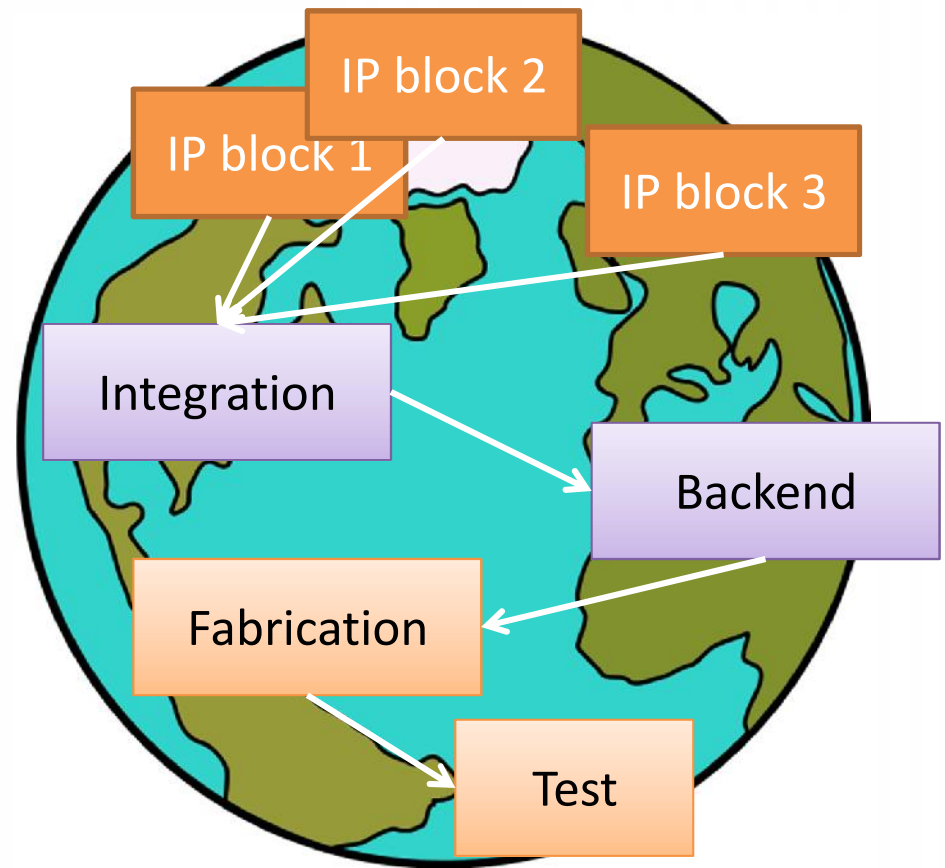
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Technion – Israel Institute of Technology
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University of Haifa and Intel Israel

Outline

- IP theft issue in SoC
- Reverse Engineering with Scan
- Junta Learning
- Clustering and Graph Completion
- The Test Case: BitCoin SHA-256
- Conclusions

IP Piracy

- Modern SoC development mode: global and distributed
- IP passes dozens of hands
- Issue of Trust



Preventing IP theft

- **Watermarks** – allow identification without altering the function
 - State Machine Encoding
 - Constraints on physical layout
 - More...
 - Detection
 - Proof
- Forensic techniques
 - Direct detection



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Reverse Engineering of an ASIC

- Phase 1 – Invasive Physical → Circuit
 - Delayering
 - SEM
 - Nanoscale Imaging
 - Cross-section
- Phase 2 – Algorithmic Circuit → Spec
 - FSM Extraction
 - Model Checking
 - SAT

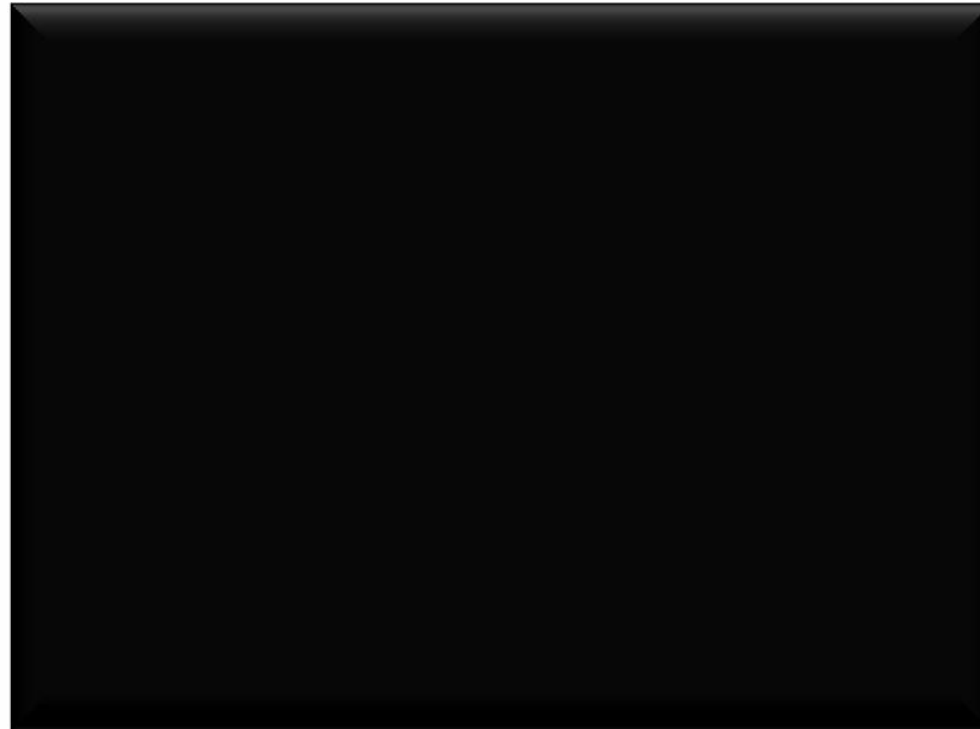


Reverse Engineering of an ASIC

- Phase 1 – Invasive Physical → Circuit
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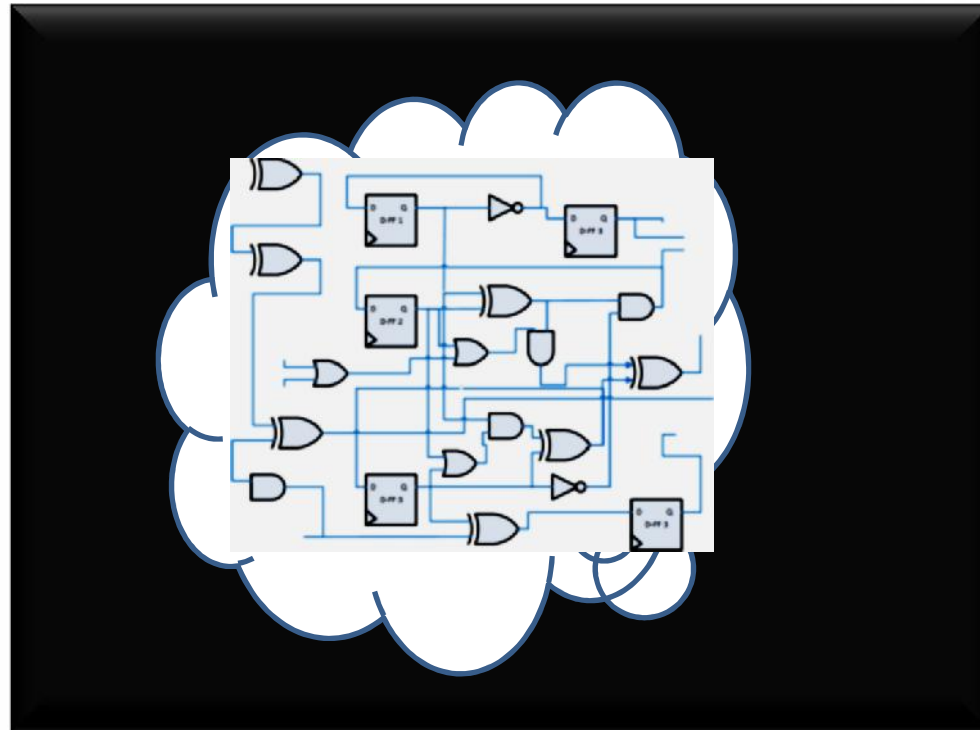
Scan Side Channel makes phase 1 non-invasive

The Scan Technique



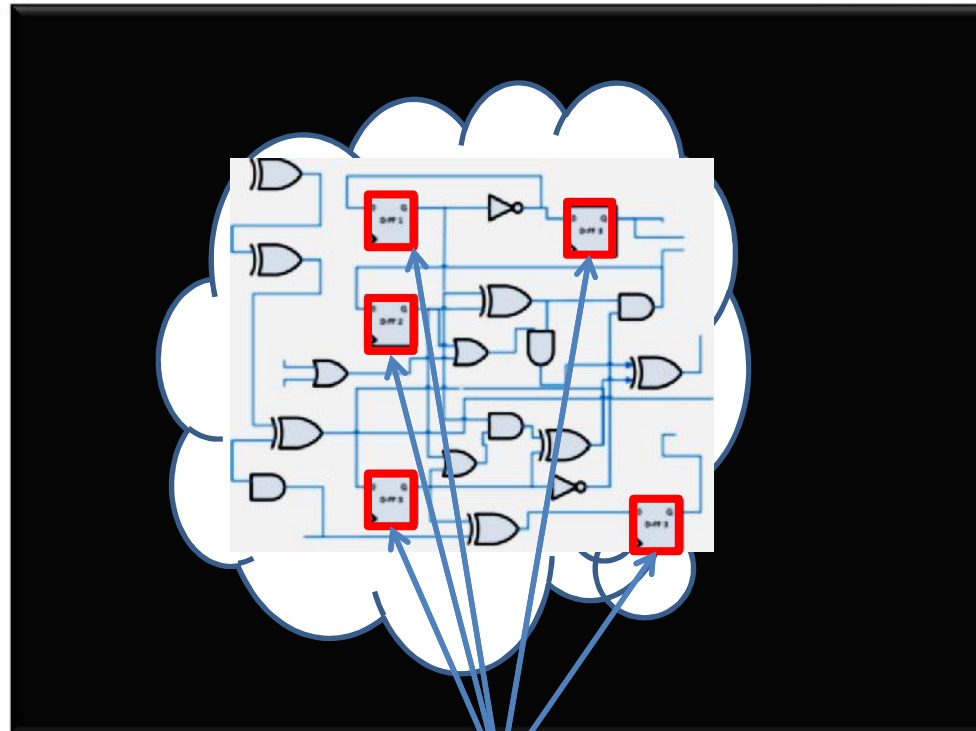
Goal: automate production testing

The Scan Technique



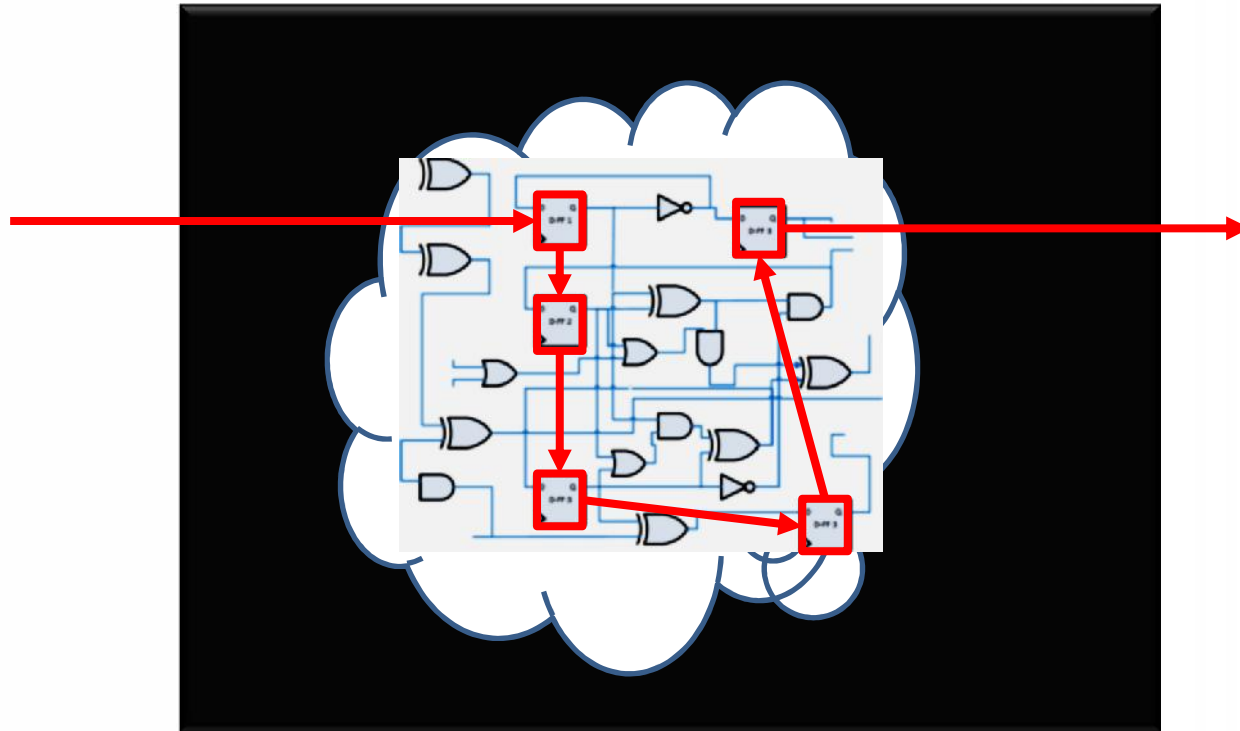
Need to verify every net is functional

The Scan Technique



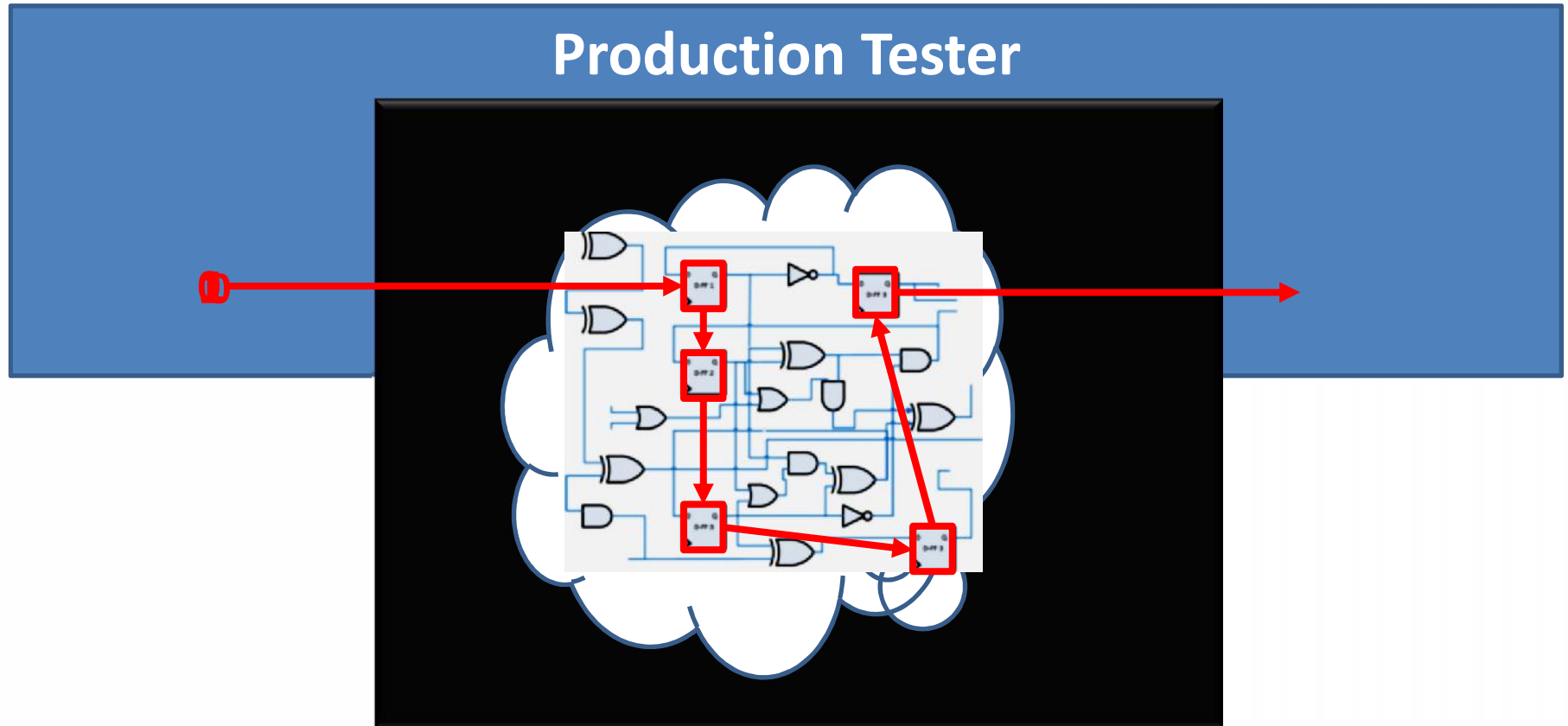
Sequential Cells
(FFs / Latches)

The Scan Technique



Scan Insertion

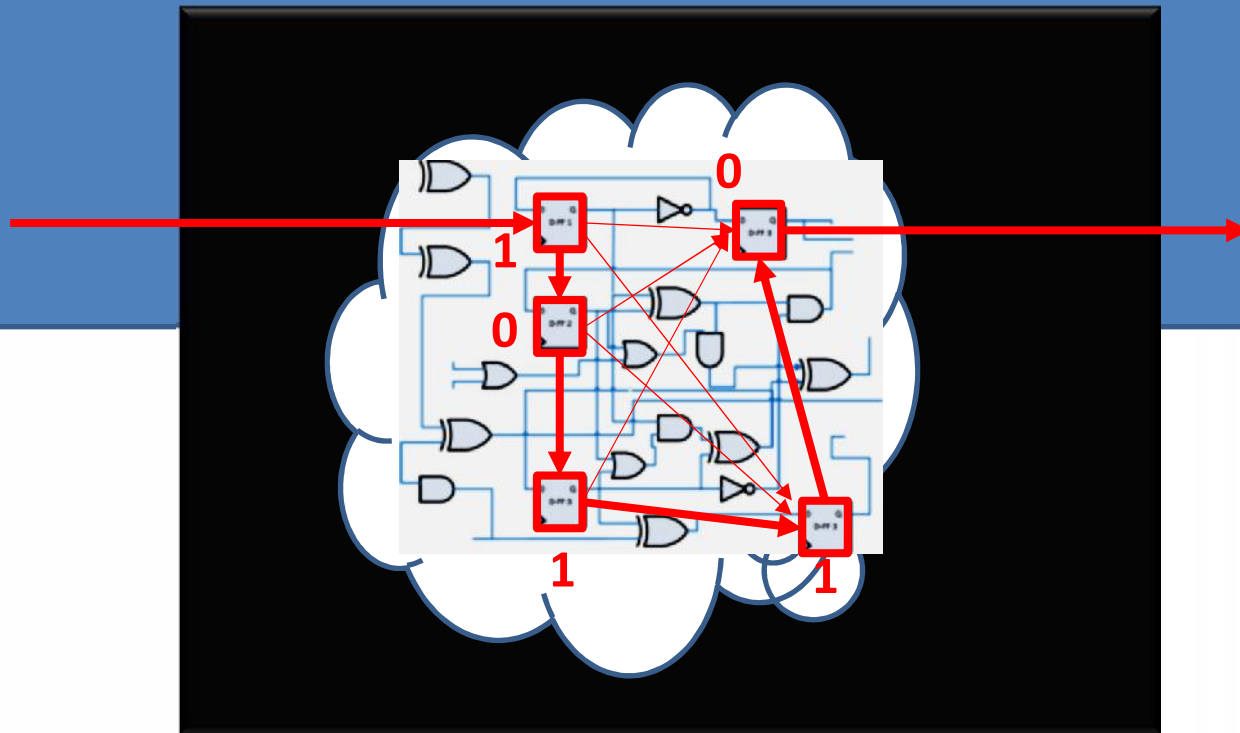
The Scan Technique



Shift In

The Scan Technique

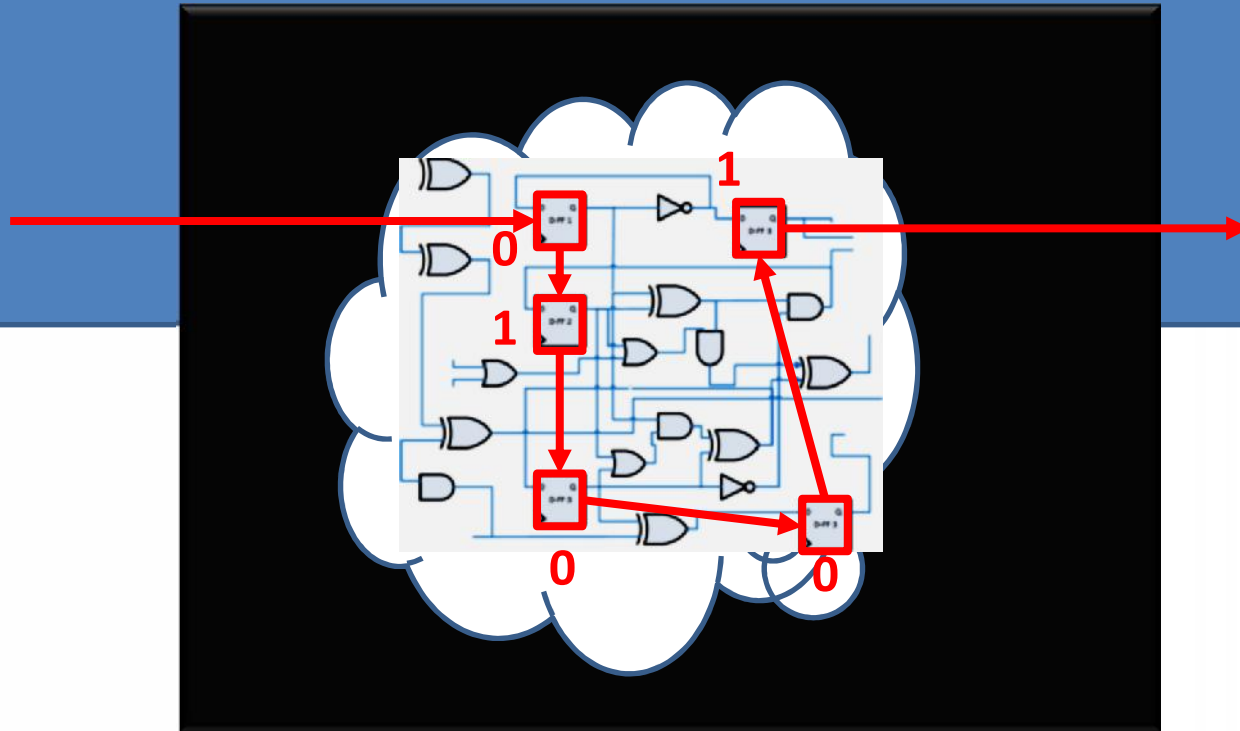
Production Tester



Capture

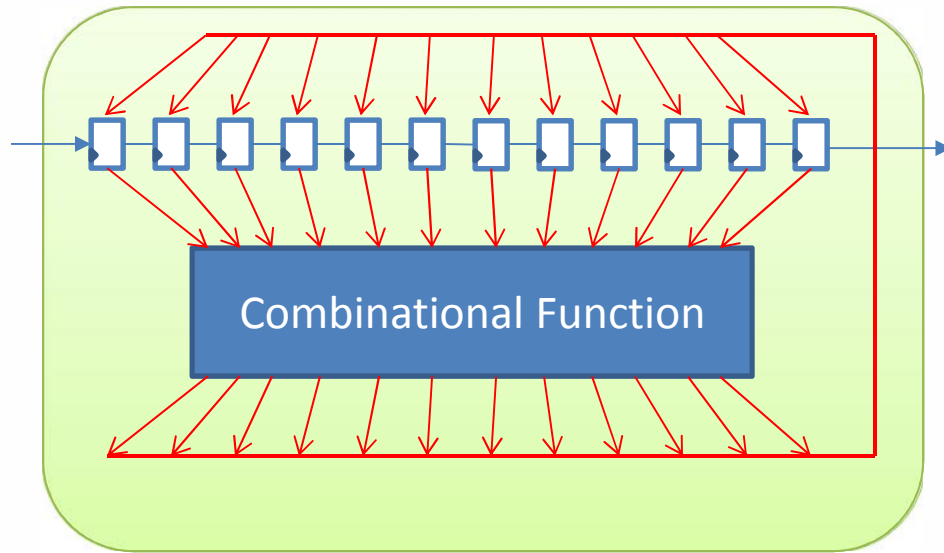
The Scan Technique

Production Tester



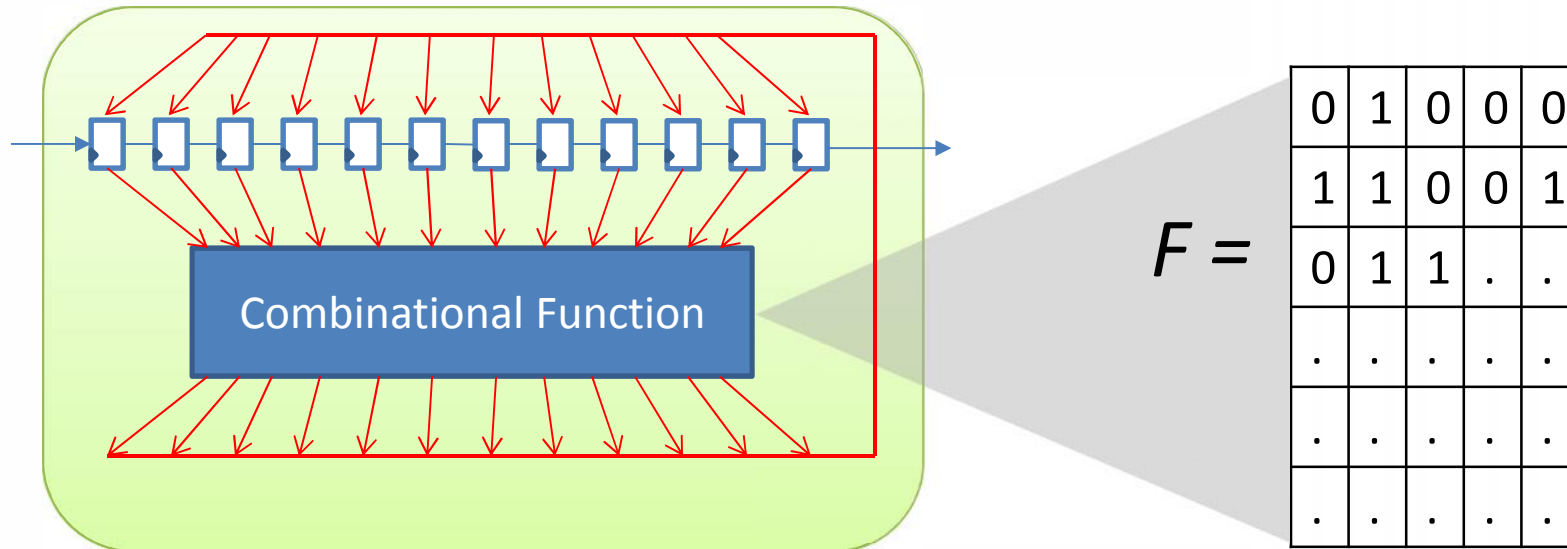
Shift Out

Unfolding Sequential Circuits with Scan



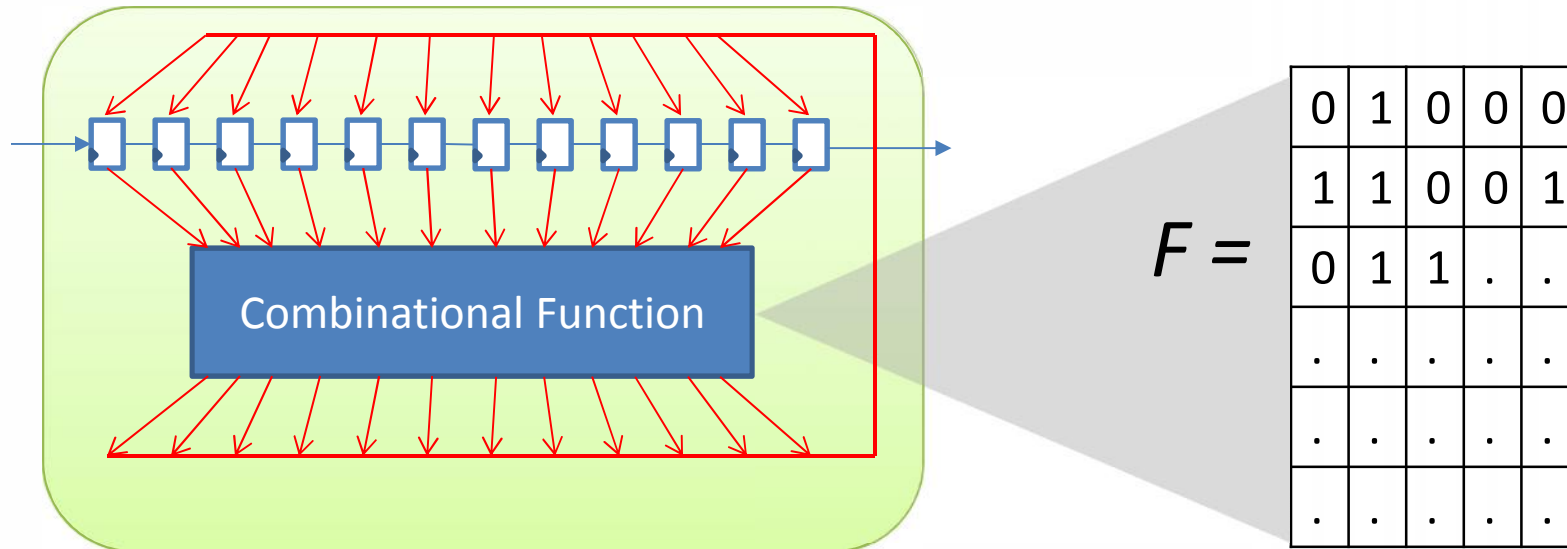
- Scan turns the SoC to a stateless circuit
- Mapped to the **Boolean Function Learning** problem: $\{0,1\}^n \rightarrow \{0,1\}^n$

Unfolding Sequential Circuits with Scan



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- Exhaustive Search: Extract the Truth Table by running queries for all inputs

Unfolding Sequential Circuits with Scan



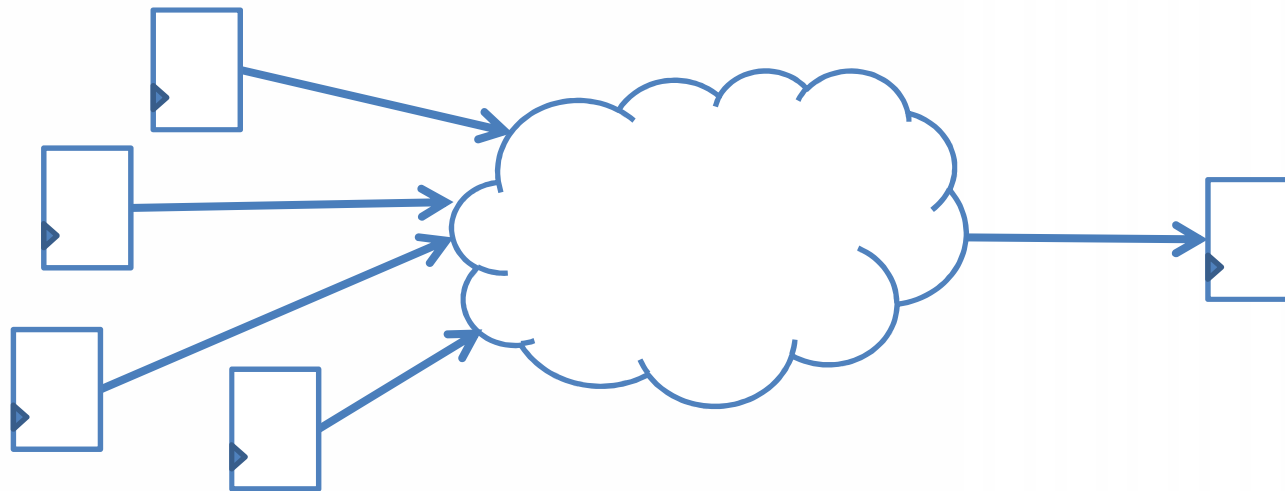
- Scan turns the ASIC to a stateless circuit
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- **Exponential Size: 2^n**

Outline

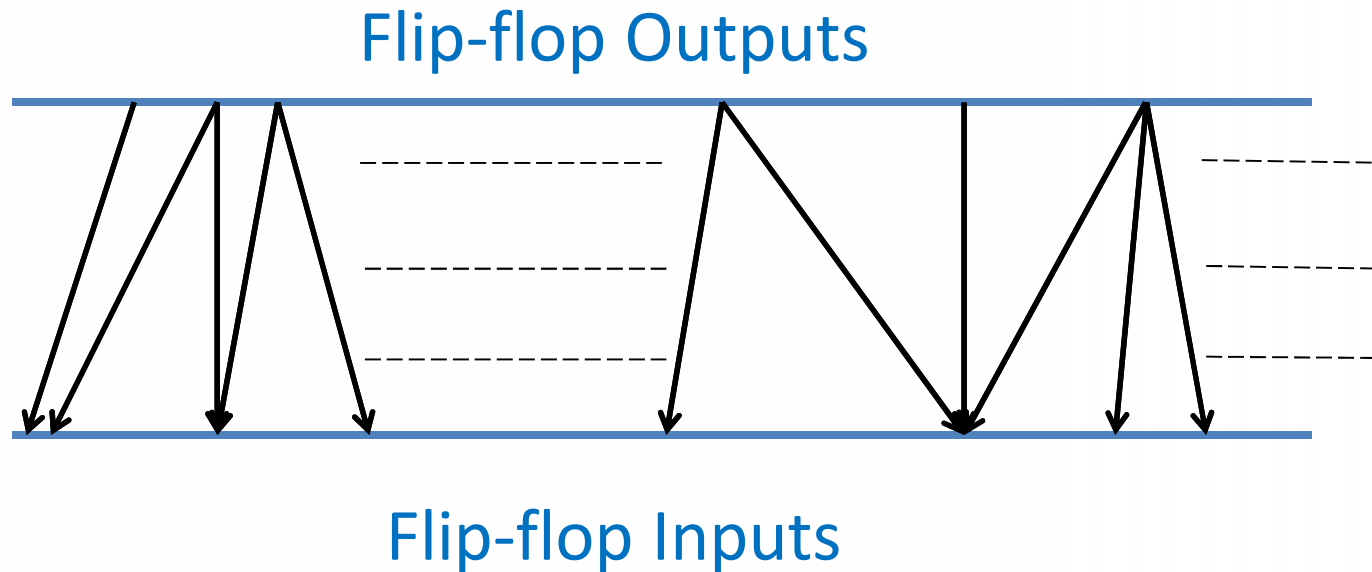
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Limited Transitive Fan-in

- In practice, logic cones have limited number of inputs: Transitive Fan In = K



Dependency Graph



- Bipartite graph represents flip-flop dependencies
- The goal: Find dependencies
- Complexity: $2^n \rightarrow 2^k$: Scalable with the chip size

The K-Junta Algorithm

$$y = f(\vec{x}), \vec{x} = \{x_1, x_2, \dots, x_i, x_{i+1}, \dots, x_j, \dots, x_n\}$$

Generate random queries $y = f(\vec{x})$

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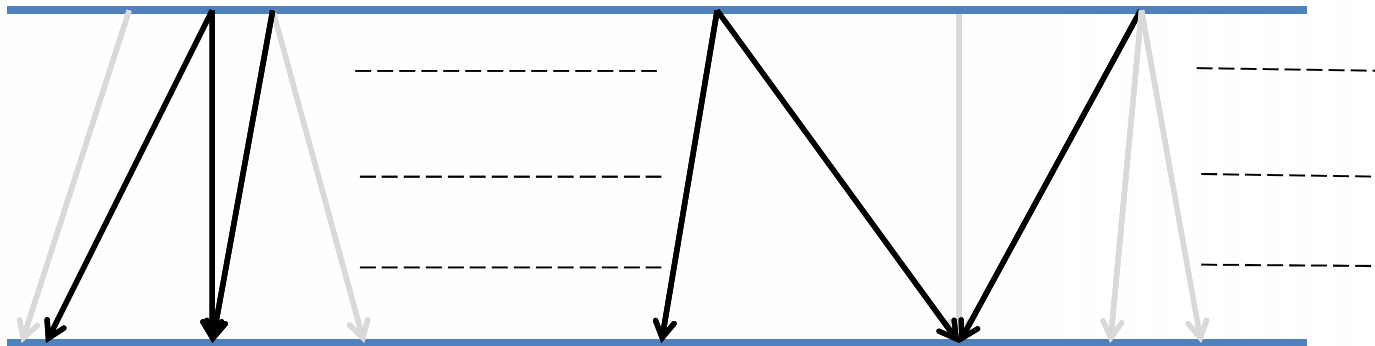
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$$O(n \cdot \log n \cdot k \cdot 2^k)$$

Relevant Variable

Partial Dependency Graph

Flip-flop Outputs



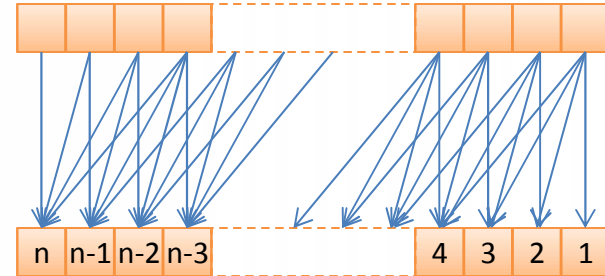
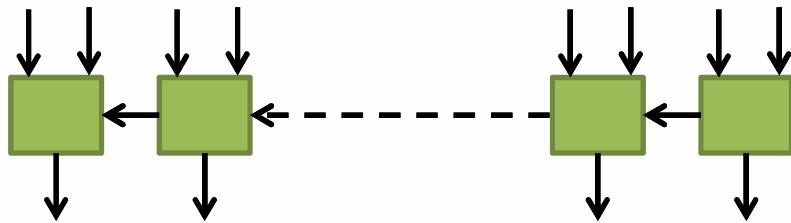
Flip-flop Inputs

- If k is too high \rightarrow Partial dependency graph
- **Influence** = sensitivity of a function to a variable
- K -Junta works for Influence $> 1/2^K$

Outline

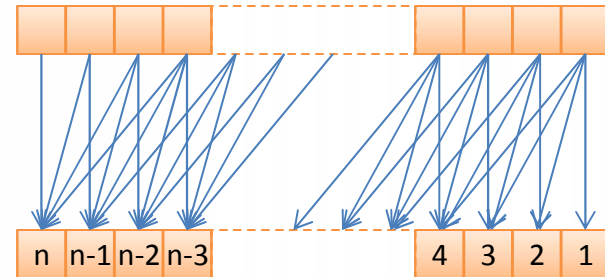
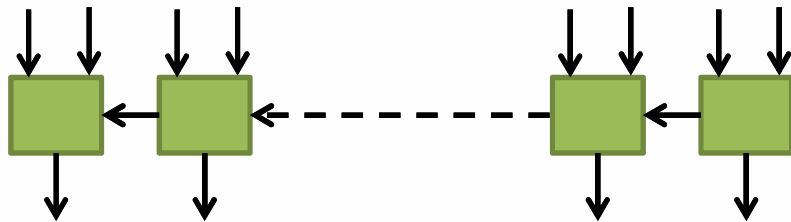
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The Adder Example



- Dependencies across many bits are not likely to appear
 - Influence too low
- Close neighbor dependencies are discovered
- Need to group all the nodes of the adder

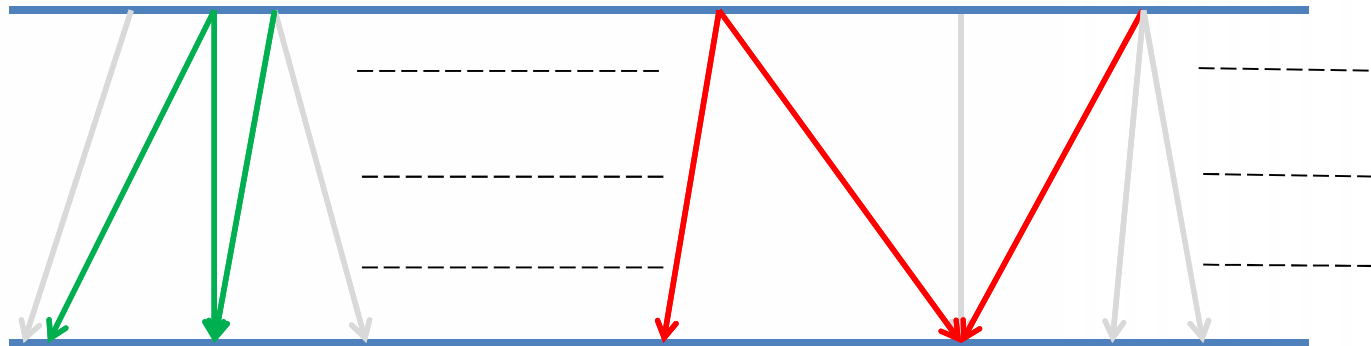
SNN Clustering



- Shared Nearest Neighbors Clustering
 - Every pair of nodes with >threshold shared dependencies assigned to the same cluster

SNN Clustering

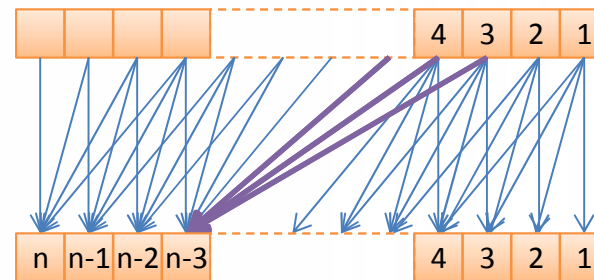
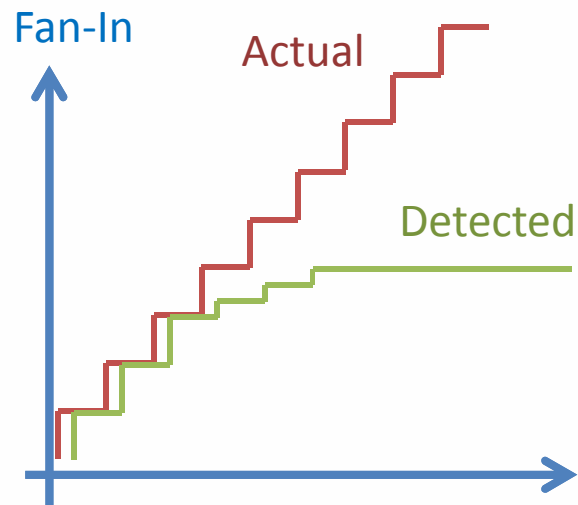
Flip-flop Outputs



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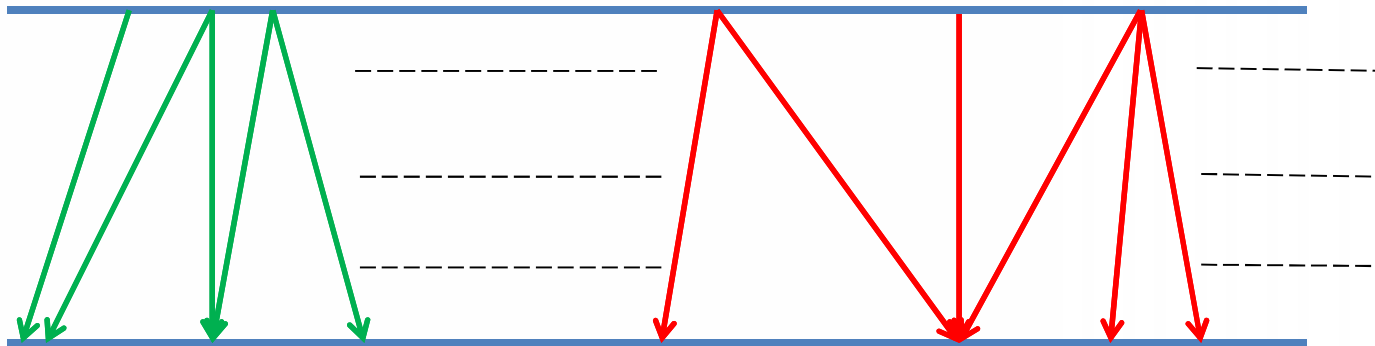
Enumeration of the Adder Nodes



- Sort outputs in a cluster by their fan-in
 - Sort inputs accordingly
- Handle the plateau by iterative enumeration
 - Higher order inputs feed higher order outputs

Completing the graph

Flip-flop Outputs



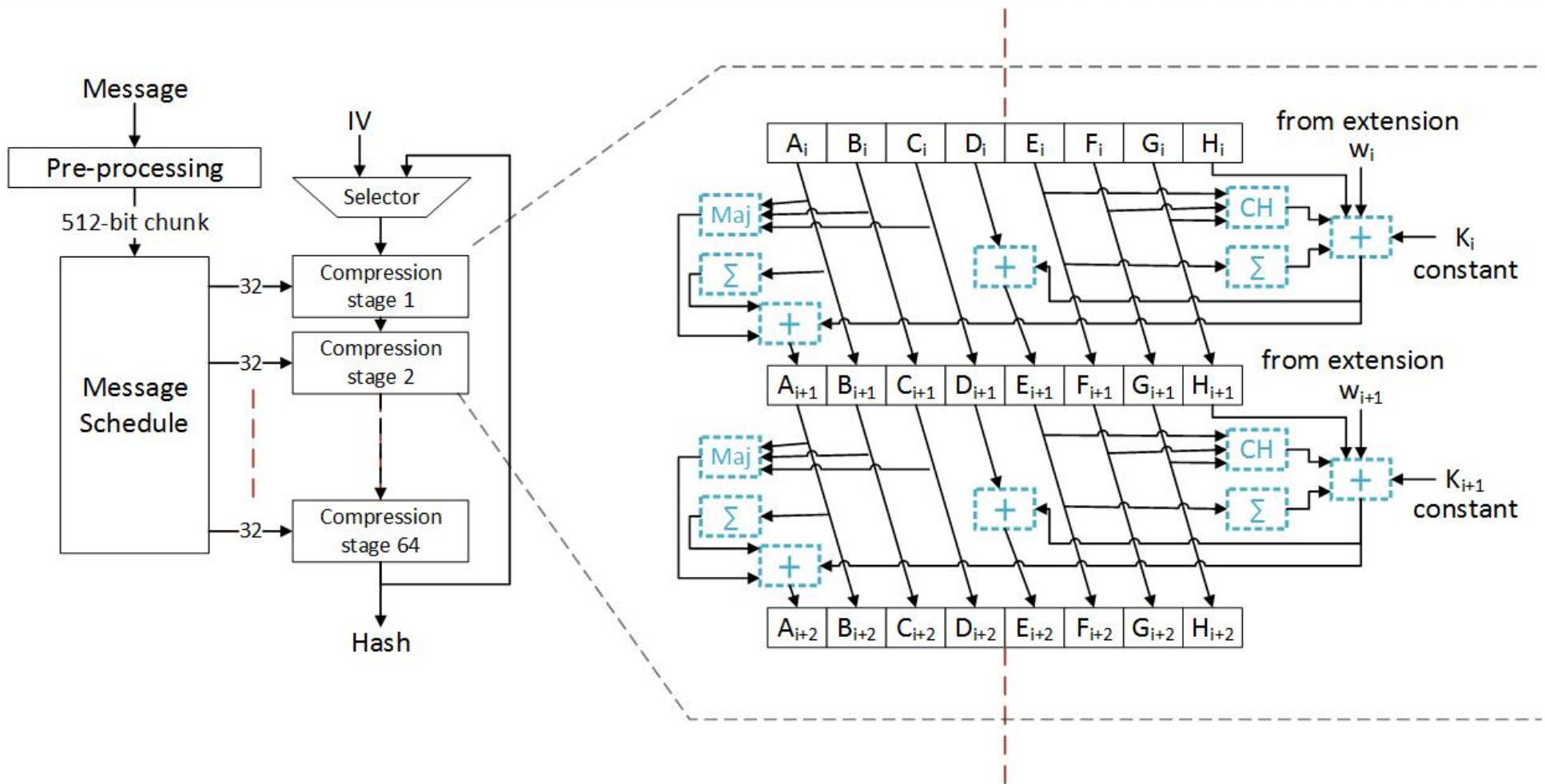
Flip-flop Inputs

- Assuming the learner is looking for an adder
- Add dependencies of output bit i on all input bits 0 to i .

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SHA-256 Structure



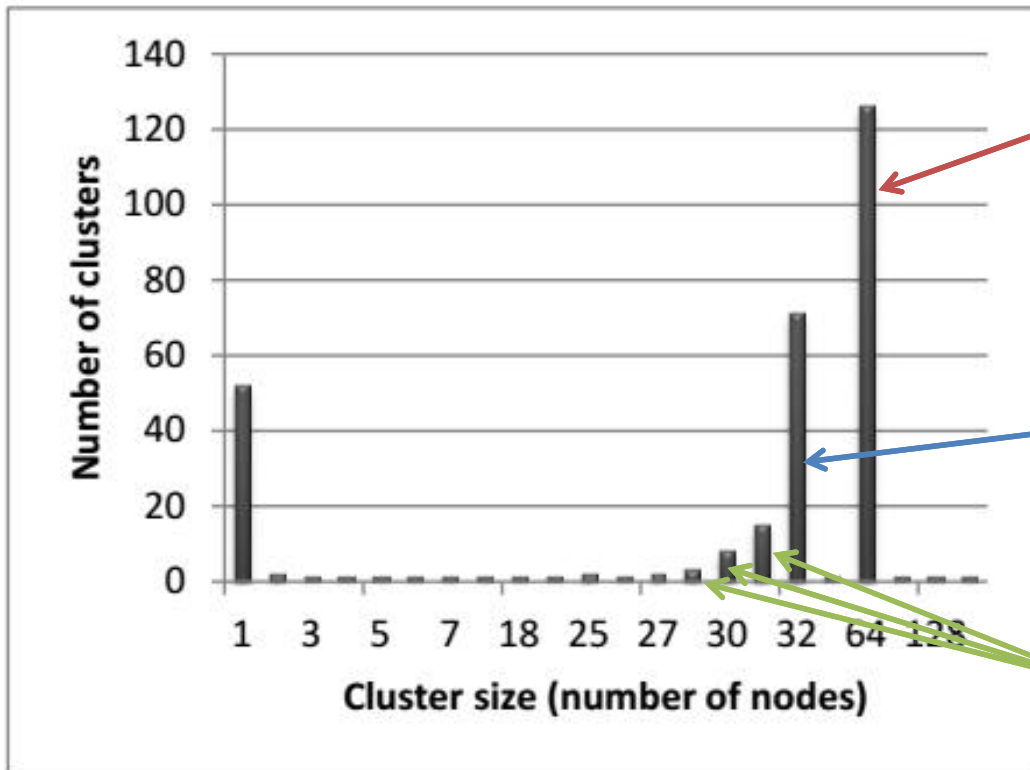
Learning Strategy

- The implementation is not known in advance
- But there are building blocks inherent to SHA-256
 - 7-way adder
 - 5-way adder
- We search for structures that look like adders

BitCoin SHA-256 Accelerator

- Open source design from opencores.org
- Performance oriented, heavily pipelined
- ~80,000 registers
- Used a software simulator

After K-Junta and Clustering



64-sized clusters match
2 32-bit adders
→ Compression Stage

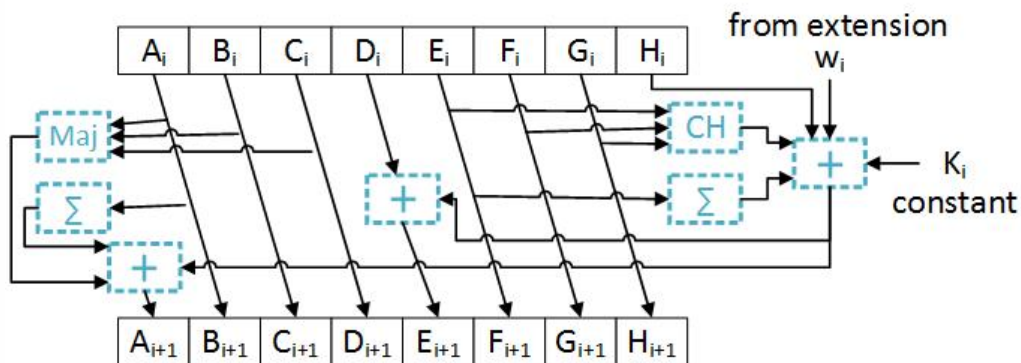
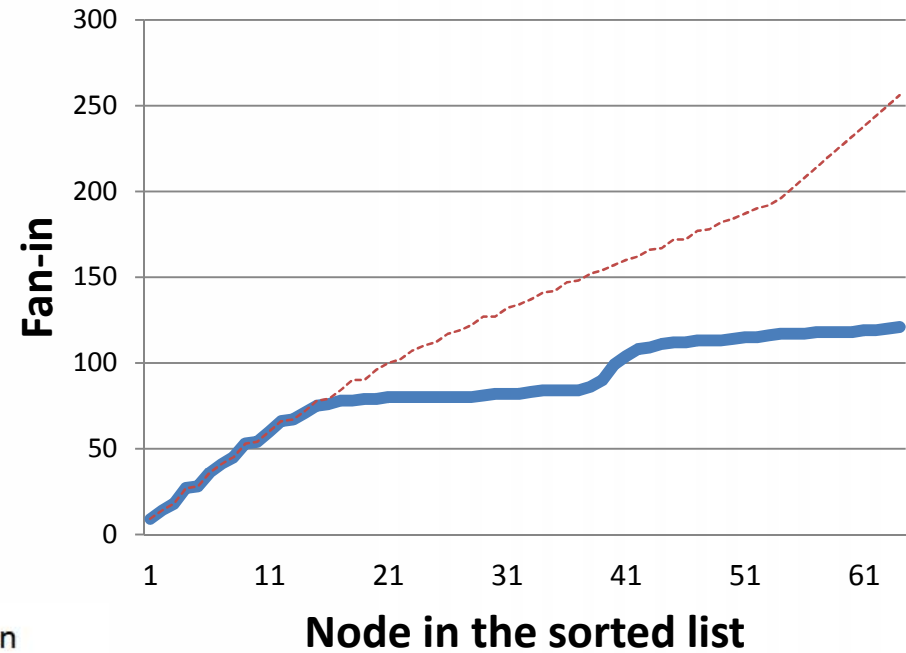
32-sized clusters match
1 32-bit adder
→ Message Schedule

SNN Clustering Error

Number of stages suggests two SHA-256 instances, but not necessarily

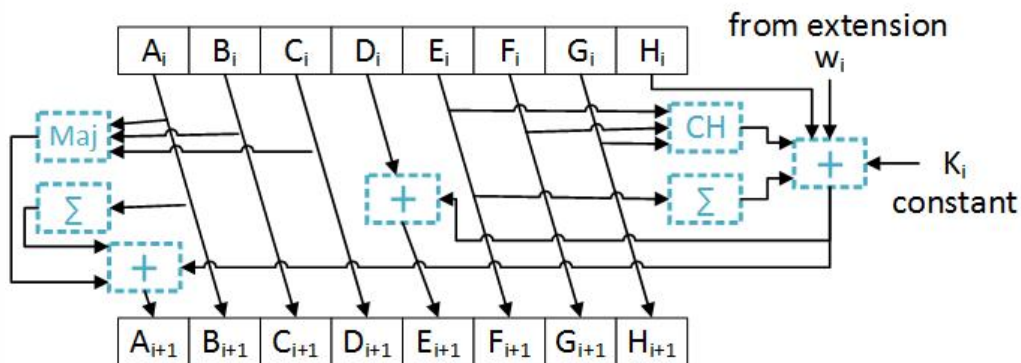
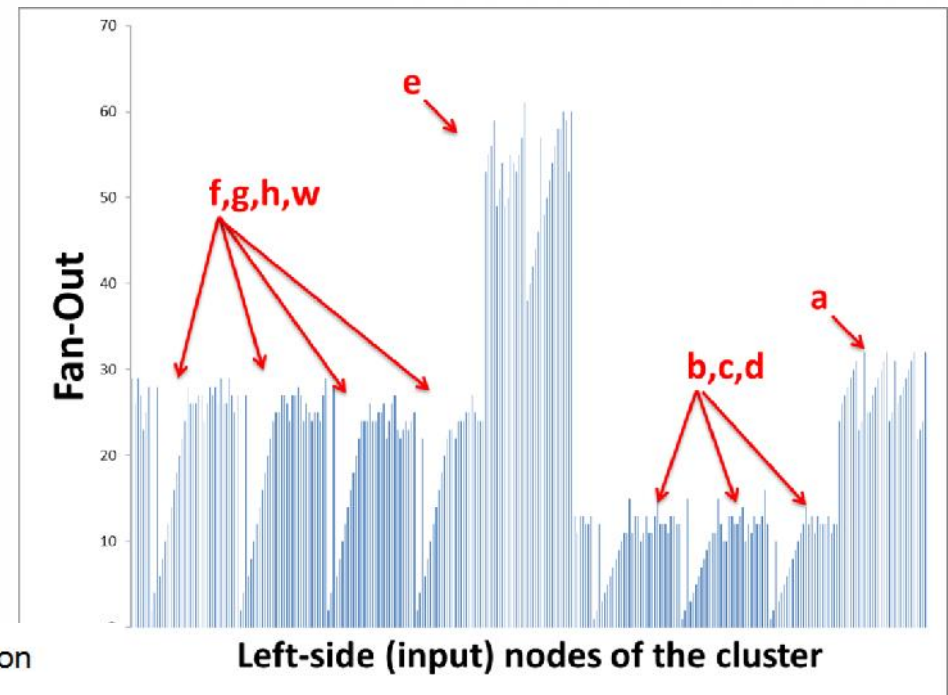
Zooming in into a cluster

- Sort by enumeration
- How to detect individual operands?

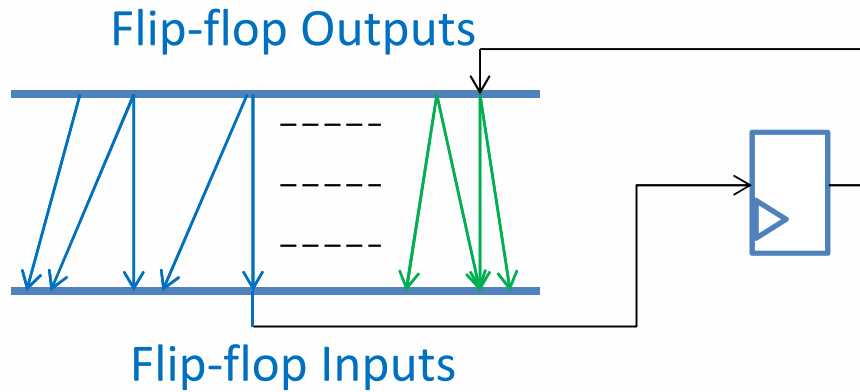


Detecting operands by fanout

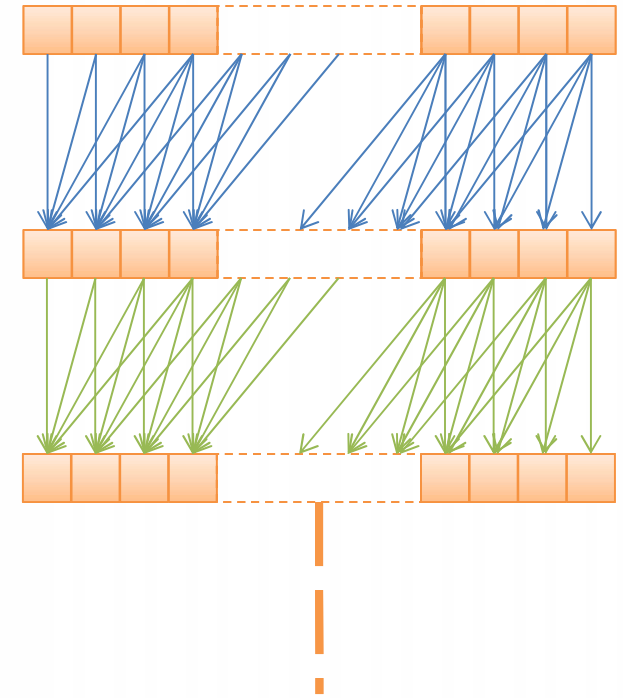
- Fanout components
 - Bit order
 - Number of functions
 - Function type



Returning to sequential



Flattened



Folded

Summary

- A novel method of IP theft detection
 - By non-invasive reverse engineering with scan
 - Boolean function analysis and graph methods
 - Works with or without watermarks
- Learned a 80,000-register SHA-256 accelerator
- What next
 - More test cases
 - Detecting Trojan hardware

Thanks!